HIGH PERFORMANCE SWITCHES AND ROUTERS Wiley H. JONATHAN CHAO and BIN LIU Instructor: Mansour Rousta Zadeh

Knockout Switches Outlines

- Introduction
- Single Stage Knockout-Basic Architecture
- Knockout Concentration Principle
- Concentrator Architecture
- Channel grouping Principle
- Generalized Knockout Principle
- MOBAS Switch
- Fault Tolerant Knockout Switches
- Conclusion

Knockout Switches Introduction-1

- Output Buffered Switches: The best delaythroughput performance.
- Problem of Output Buffered Switches:
 - Memory Speed Limitation
- Solution (Knockout Principle)
 - limiting the number of cells that can arrive at an output port in each time slot
 - Other cells are discarded
- **Q**:How Many?
- Tradeoff between "Cell Loss Ratio" and "Memory Bandwidth"

Introduction-2

- the memory speed is no longer the bottleneck for the output-buffered switch
- No Commercial Products

Why?

Single Stage Knockout-Basic Architecture



Knockout switch interconnection fabric

Knockout Switches-Bus Interface



Operation of a Barrel Shifter-1



Operation of a Barrel Shifter-2



Knockout Concentration Principle

$$P_k = \binom{N}{k} \left(\frac{\rho}{N}\right)^k \left(1 - \frac{\rho}{N}\right)^{N-k}, \qquad k = 0, 1, \dots, N.$$

$$\Pr[\text{cell loss}] = \frac{1}{\rho} \sum_{k=L+1}^{N} (k-L) {\binom{N}{k}} \cdot \left(\frac{\rho}{N}\right)^{k} \left(1 - \frac{\rho}{N}\right)^{N-k}$$

$$\Pr[\text{cell loss}] = \left(1 - \frac{L}{\rho}\right) \left(1 - \sum_{k=0}^{L} \frac{\rho^k e^{-\rho}}{k!}\right) + \frac{\rho^L e^{-\rho}}{L!} \qquad N \to \infty,$$

Concentration Cell Loss Performance



Knockout Switches Concentration Cell Loss Performance



Construction of the Concentrator



Winner Loser

(a)

(b)

An eight-input to four-output concentrator



Outputs

Construction of large Concentrator with Small Concentrators



Knockout Switches channel grouping principle



An asymmetric switch with line expansion ratio KM/N



Knockout Switches Maximum Throughput

TABLE 6.1 Maximum Throughput with $K \neq N$ Kept Constant While $K, N \rightarrow \infty$

	Maximum Throughput									
М	K/N =	$\frac{1}{16}$	$\frac{1}{8}$	$\frac{1}{4}$	$\frac{1}{2}$	1	2	4	8	16
1 2 4 8 16		0.061 0.121 0.241 0.476 0.878	0.117 0.233 0.457 0.831 0.999	0.219 0.426 0.768 0.991 1.000	0.382 0.686 0.959 1.000	0.586 0.885 0.996 1.000	0.764 0.966 1.000	0.877 0.991 1.000	0.938 0.998 1.000	0.969 0.999

	Maximum Throughput									
M	KM/N =	1	2	4	8	16	32			
1		0.586	0.764	0.877	0.938	0.969	0.984			
2		0.686	0.885	0.966	0.991	0.998	0.999			
4		0.768	0.959	0.996	1.000	1.000	1.000			
8		0.831	0.991	1.000						
16		0.878	0.999							
32		0.912	1.000							
64		0.937								
128		0.955								
256		0.968								
512		0.978								
1024		0.984								

TABLE 6.2 Maximum Throughput with KM / N Kept Constant While KM, $N \rightarrow \infty$

Generalized Knockout Principle



Generalized Knockout Principle

$$\Pr[\text{cell loss}] = \frac{1}{M\rho} \sum_{k=LM+1}^{N} (k - LM) {\binom{N}{k}} {\left(\frac{M\rho}{N}\right)^{k}} \cdot \left(1 - \frac{M\rho}{N}\right)^{N-k}$$

$$\Pr[\text{cell loss}] = \left(1 - \frac{L}{\rho}\right) \left(1 - \sum_{k=0}^{LM} \frac{(M\rho)^k e^{-M\rho}}{k!}\right) + \frac{(M\rho)^{LM} e^{-M\rho}}{(LM)!}$$

$$\text{As } N \to \infty$$

Generalized knockout principle Operation



number of simultaneous cells accepted, $L \times M$



The architecture of a multicast output buffered ATM switch



Replicating cells for a multicast connection in the MOBAS





OPC : Output Port Controller

Knockout Switches Multicast Grouping Network



Switching condition of the switch element SWE



Knockout Switches Routing a multicast cell



Knockout Switches Translation Tables-1

(a) Unicast translation table entry at IPC1	Arrived VCI	A1 A2		Р	New VCI	
(b) Multicast translation table entry at IPC1	Arrived VCI	MP1		Р	BCN	
(c) Multicast translation table entry at IPC2	BCN	MP2				
(d) Multicast translation table entry at OPC	BCN	DR	VCI	1	VCI 2	

- VCI : Virtual channel identifier
- A1 : Output address of MGN1
- MP1 : Multicast pattern in MGN1
- P : Priority for cell contention

- BCN : Broadcast channel number
- A2 : Output address of MGN2
- MP2 : Multicast pattern in MGN2
- DR : Number of duplication requests

Knockout Switches Translation Tables-2



I : Multicast indication bit, 0 : unicast, 1 : multicast

(b)



(a)

Fault Model of Switching Elements-1

- Need to Reliability
- Redundancy Techniques:
 - Time Redundancy
 - Space Redundancy
- Fault Tolerant Topics
 - Fault Diagnosis
 - Fault Detection
 - Fault Location
 - System Reconfiguration

Fault Model of Switching Elements-2

Fault Sources

In Control Logic Circuit (logic errors)
 Cross Stuck (Remains in Cross state)
 Toggle Stuck (Remains in Toggle state)

- In Data Link of SWE (link errors)
 - Vertical Stuck
 - Horizontal Stuck

Knockout Switches Cross-Stuck (CS) Fault



Toggle-Stuck (TS) Fault



Knockout Switches Vertical/Horizontal-Stuck(VS/HS) Fault



Knockout Switches Fault Detection

- Using FD's (Fault detector)
- Test Pattern Generation by MPM's and AB's
- On Fault Diagnosis:
 - Keeping user packets
 - Doing Tests
 - System Reconfiguration
- Fast, So no interruption in Switch Functionality

Cross-Stuck and Toggle-Stuck Fault Detection

- Toggle/Cross Fault Detection?
 - Monitoring FA's
 - If all 0 -> Fault
- Online Detection?
 - Toggle Stuck –possible!
 - Cross Stuck –Not possible!
- Vertial/Hor. Stuck Fault Detection?
 - Monitoring Switch module
 If all the same ->Fault

Fault Location and Reconfiguration

- Fault localization after Detection
- System Reconfiguration after localization
- Using test cells
 - FileIds:
 - FA Field
 - New Priority field $[\log_2 2L_1 M]$ bits in MGN1
 - Input source field

Toggle-Stuck and Cross-Stuck Cases



Knockout Switches Fault Location Test



Fault Location of Cross Stuck



Reconfiguration of Vertical Stuck Fault



Knockout Switches Fault Location Test



Knockout Switches Reconfiguration

